

# Category Theory Exercise Sheet 1

September 10, 2013

The deadline is 6pm on Thursday the 19th of September. You can either email your answers to `r.furber` at `cs.ru.nl` or put them in the Robert Furber's postbox on the north side of the second floor of the Huygensgebouw. It is inside a white cabinet with its back to the stairs, opposite the photocopier.

For each question, the number of marks available is indicated in round brackets. The total number is 50.

## 1

In this question, you will prove some basic facts about isomorphisms that are valid in any category. When proving two expressions are equal, give full details using the axiomatic definition of a category.

(a) Let  $f : X \rightarrow Y$  be a map. If  $g, g' : Y \rightarrow X$  are both inverses of  $f$ , then  $g = g'$ .

(3)

(b) Show that  $\text{id}_X : X \rightarrow X$  is an isomorphism.

(2)

(c) Show that if  $f : X \rightarrow Y, g : Y \rightarrow Z$  are isomorphisms,  $g \circ f$  is an isomorphism.

(5)

(d)  $\text{Aut}(X)$  is the set of isomorphisms  $X \rightarrow X$ . What algebraic structure does it have?

(2)

## 2 (Awodey 1.9.3)

Show that

- (a) In **Sets**, the isomorphisms are exactly the bijections.

(7)

- (b) In **Monoids**, the isomorphisms are exactly the bijective homomorphisms.

(4)

- (c) In **Posets**, the isomorphisms are not the same as the bijective homomorphisms.

(6)

## 3

- (a) Show that a monoid homomorphism  $f : \mathbb{N} \rightarrow \mathbb{Z}$  is determined by  $f(1)$ .

(2)

- (b) Show that such an  $f$  can never have an inverse (in **Monoids**).

(4)

## 4 [1, page 9]

The arrows in a category form a structure in their own right. We can define it as follows.

Consider a set  $A$ , with a relation  $R \subseteq A \times A \times A$ . We write the relation  $R(f, g, h)$  as  $f \bullet g = h$ , and we say that  $(f, g)$  is *composable* iff there exists an  $h$  such that  $f \bullet g = h$ . An identity element of  $A$  is a  $u \in A$  such that  $f \bullet u = f$  whenever  $(f, u)$  is composable and  $u \bullet g = g$  whenever  $(u, g)$  is composable.

We define an *arrows-only-category* as a pair  $(A, R \subseteq A \times A \times A)$  such that:

- (i) If  $f \bullet g = h$  and  $f \bullet g = h'$  then  $h = h'$ .
- (ii)  $(k \bullet g) \bullet f$  is defined iff  $k \bullet (g \bullet f)$  is defined, and  $(k \bullet g) \bullet f = k \bullet (g \bullet f)$  in this case.
- (iii) If  $(k, g)$  and  $(g, f)$  are composable, then  $(k \bullet g, f)$  (equivalently  $(k, g \bullet f)$ ) is composable.
- (iv) For each  $g \in A$ , there exist identity arrows  $u, u' \in A$  such that  $(u, g)$  and  $(g, u')$  are composable.

## 4.1 Questions

- (a) Show that the arrows of every category make an arrows-only-category. (5)
- (b) Show that the identity arrows in an arrows-only-category are unique, and that an identity arrow is equal to its own identity arrows. (5)
- (c) Show that from every arrows-only-category one can recover the objects to make a category. (5)

## References

- [1] Saunders Mac Lane. *Categories for the Working Mathematician*. Graduate Texts in Mathematics. Springer Verlag, 1971.